

FACULTY OF ENGINEERING & TECHNOLOGY

BCA-307 Operating System

Lecturer-09

Manisha Verma

Assistant Professor Computer Science & Engineering

Process Synchronization

Background The Critical-Section Problem



Background

•Processes can execute concurrently

May be interrupted at any time, partially completing execution

•Concurrent access to shared data may result in data inconsistency.

•Maintaining data consistency requires mechanisms to ensure the orderly execution of cooperating processes

•Illustration of the problem:-

Suppose that we wanted to provide a solution to the consumer-producer problem that fills **all** the buffers. We can do so by having an integer **counter** that keeps track of the number of full buffers. Initially, **counter** is set to 0. It is incremented by the producer after it produces a new buffer and is decremented by the consumer after it consumes a buffer.

Producer

while (true) {

/* produce an item in next produced */

while (counter == BUFFER_SIZE);

/* do nothing */

buffer[in] = next_produced;

in = (in + 1) % BUFFER_SIZE;

counter++;



Consumer

```
while (true)
{
    while (counter == 0)
        ; /* do nothing */
        next_consumed = buffer[out];
        out = (out + 1) % BUFFER_SIZE;
        counter--;
        /* consume the item in next consumed */
}
```

Critical Section Problem

•Consider system of *n* processes $\{p_0, p_1, \dots, p_{n-1}\}$

•Each process has critical section segment of code

Process may be changing common variables, updating table, writing file, etc
 When one process in critical section, no other may be in its critical section

•Critical section problem is to design protocol to solve this...

>Each process must ask permission to enter critical section in entry section, may follow critical section with exit section, then remainder section



structure of process P_i

do {

entry section

critical section

exit section

remainder section

} while (true);

1. Mutual Exclusion - If process P_i is executing in its critical section, then no other processes can be executing in their critical sections

- 2. Progress If no process is executing in its critical section and there exist some processes that wish to enter their critical section, then the selection of the processes that will enter the critical section next cannot be postponed indefinitely
- 3. Bounded Waiting A bound must exist on the number of times that other processes are allowed to enter their critical sections after a process has made a request to enter its critical section and before that request is granted
 - Assume that each process executes at a nonzero speed
 - No assumption concerning relative speed of the *n* processes

Critical-Section Handling in OS:-----

Two approaches depending on if kernel is preemptive or non-preemptive

- Preemptive allows preemption of process when running in kernel mode
- Non-preemptive runs until exits kernel mode, blocks, or voluntarily yields CPU
 - Essentially free of race conditions in kernel mode

MCQ

Mutual Exclusion enter...

- A. Critical section single process
- B. Critical section no process
- C. Critical section atleast two process
- A. None

Mutual Exclusion is allowed ..

- A. Interleaving
- B. Sharable
- C. Mutable lock
- D. None

Preemptive – allows preemption of process when.....

- A. running in kernel mode
- B. running in user mode
- C. running in compiler mode
- D. None



Concurrent access to.....

- A. shared data
- B. Not sharable data
- C. Mutable data
- D. None

Kernel is a....

- A. Os
- B. Compiler
- C. Process
- D. CPU

